§6 Randomization: Motivation

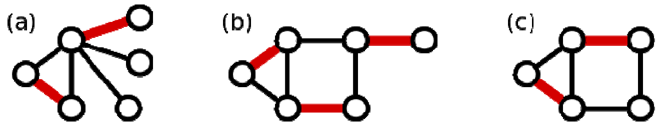
Simple polyn.-time decision whether a (not necessarily bipartite nor planar) graph admits a perfect matching.

Let x_{ij}, $1 \leq i < j \leq n$, denote variables and consider Tutte's skew-symmetric *symbolic* matrix A_G with entries

$a_{ij} := x_{ij}$ if $\{i,j\} \in E$ and $i<j$

$a_{ij} := -x_{ji}$ if $\{i,j\} \in E$ and $i>j$

$a_{ij} := 0$ otherwise.



$\det(A_G) = \sum_\pi \mathrm{sign}(\pi) \cdot a_{1,\pi(1)} \cdot a_{2,\pi(2)} \cdot a_{3,\pi(3)} \cdots a_{n,\pi(n)}$

- is an n^2-variate integer polynomial of total degree n
- that can be evaluated using $O(n^3)$ tests & arith. op.s
- is identically zero iff G has *no* perfect matching!

Recall: *A perfect matching* in a graph $G=(V,E)$ of $|V|=2n$ vertices is a set $M \subseteq E$ of n edges without common vertices.

Lemma on Tutte's Determinant

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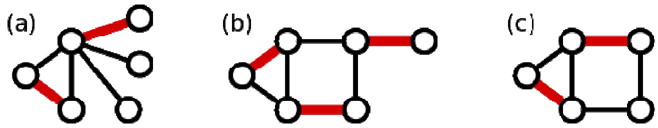

Proof '$\Rightarrow$' A perfect matching is a permutation $\mu: V \to V$ s.t. $\forall i$: $\{i,\mu(i)\} \in E$ (*) and all cycles have length 2. Set $x_{i,\mu(i)} := 1$, $x_{ij} := 0$ for $j \neq \mu(i)$. Then $\det(A_G)(\underline{x}) = 1$ (why?)

'$\Leftarrow$' Let $\det(A_G) = \sum'_{\pi \text{ has odd cycle}} + \sum''_{\pi \text{ only of even cycles}}$

Then $\sum'_\pi \equiv 0$. Let π consist of only even cycles s.t. (*). This gives rise to a perfect matching.

Recap: symmetry, cycle decompos., multivar. polyn.

Polynomial Identity Testing

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$\det(A_G) \;=\; \sum_\pi \mathrm{sign}(\pi) \cdot a_{1,\pi(1)} \cdot a_{2,\pi(2)} \cdot a_{3,\pi(3)} \cdots a_{n,\pi(n)}$

- is identically zero iff G has *no* perfect matching;
- is an n^2-variate integer polynomial of total degree n
- that can be evaluated using $O(n^3)$ tests & arith. op.s

Recap (by example): The *total degree* of $x^2 \cdot y^3$ is 5.
Univariate polynom. of degree d has (at most) d roots.

Lemma (*Schwartz-Zippel*)**:** Fix domain D, finite $S \subseteq D$, and let $0 \neq p \in D[x_1,\ldots x_n]$ have total degree $\leq d$. Sample $r_1,\ldots r_n$ from S independently uniformly at random (*iid*).
Then (*) $\Pr[\, p(r_1,\ldots r_n)=0 \,] \leq d/|S|$.

Let j max s.t. $p_j \neq 0$

Proof (induct): $0 \neq p(x_1,\ldots x_n) = \sum_{0\leq j\leq d} p_j(x_1,\ldots x_{n-1}) \cdot x_n^{\,j}$

$(*) \leq \Pr[\, p_j(r_1,\ldots r_{n-1})=0] + \Pr[\, p(r_1,\ldots r_n)=0 \mid p_j(r_1,\ldots r_{n-1}) \neq 0]$

Markov Chain Algorithm for 3SAT

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- 1-sided error: Suppose $\underline{z}$ is a satisfying assignment
- and $\underline{y}$ guessed in line 3 differs from $\underline{z}$ at $\leq k$ places.
- After one iteration of innermost loop (lines 5 to 8):
- With probability ≥⅓ differs $\underline{y}$ only at $\leq k$-1 places.
- Loop arrives at $\underline{y}=\underline{z}$ with probability $\geq (⅓)^k$.
- Naïve choice $k{:=}n/2$ and $K{:=}3^k$.
- Better $k{:=}n/4$ and $K := 3^k \cdot 2^n / \binom{n}{k} \approx (1.5)^n$
- Current record $k{:=}3n$ and $K := (4/3^n)$

runtime $(1.33)^n \cdot \mathrm{poly}(n)$

1 Given 3CNF term $\varphi(x_1,\ldots,x_n)$
2 Repeat K times:
3 Guess assignment $\underline{y} \in \{0,1\}^n$
4 Repeat k times:
5 If $\varphi(\underline{y})=1$, accept and stop.
6 C be 1st clause in φ st $C(\underline{y})=0$
7 Guess a literal in C (1 of 3),
8 flip its assigned value in $\underline{y}$.
9 Reject!

$1/\binom{n}{cn} \approx c^{cn} \cdot (1-c)^{(1-c)n}$