

**CS500**

Spring 2016, Assignment #4

**PROBLEM 12 (6+4P) :**

a) Which of the following structures satisfy the first-order formula “ $\forall x \exists y. x + y = 0$ ”, and why not?

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|--|---|
| i) $(\mathbb{N}, +, 0)$                        | vii) $(\mathbb{Q} \setminus \{0\}, \cdot, 1)$   |
| ii) $(\mathbb{Z}, +, 0)$                       | viii) $(\mathbb{C} \setminus \{0\}, \cdot, 1)$  |
| iii) $(\{0, 1\}^d, \text{XOR}, (0, \dots, 0))$ | ix) $(\mathbb{R}^3 \setminus \{0\}, \times, (1, 1, 1))$   |
| iv) $(\mathbb{Z}, \cdot, 0)$                   | x) $(\mathbb{Q}^{2 \times 2} \setminus \{0\}, \cdot, \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix})$   |
| v) $(\mathbb{Z}, \cdot, 1)$                    | xi) $(\{ \begin{pmatrix} a & b \\ c & d \end{pmatrix} : a, b, c, d \in \mathbb{Q}, ad \neq bc \}, \cdot, \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix})$ |
| vi) $(\mathbb{Q}, \cdot, 1)$                   | xii) $(\{f : X \rightarrow X \text{ injective}\}, \circ, \text{id}), \text{ for which sets } X ?$   |

b) Consider a Boolean term  $\varphi \equiv C_1 \wedge C_2 \wedge \dots \wedge C_m$  in CNF over variables  $x_1, \dots, x_n$ .

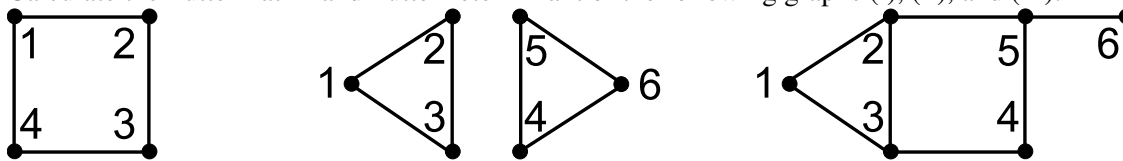
Encode it as the logical structure  $(U, V, N, P)$  with

- universe  $U = \{0, 1, \dots, n + m - 1\}$  (representing variables and clauses),
- unary relation  $V = \{0, 1, \dots, n - 1\}$  (representing variables),
- binary relation  $N = \{(v, c) : \text{variable } v \text{ occurs negated in clause } c\}$ , and
- binary relation  $P = \{(v, c) : \text{variable } v \text{ occurs unnegated in clause } c\}$ .

Moreover consider a subset  $T \subseteq U$  as an assignment of value **true** to variables  $V \cap T$  and value **false** to variables  $V \setminus T$ . Now construct a first-order formula  $\Phi(T, V, N, P)$  expressing whether  $T$  constitutes a satisfying assignment to  $\varphi$ .

**PROBLEM 13 (6+4P) :**

a) Calculate the Tutte Matrix and Tutte Determinant of the following graphs (i), (ii), and (iii):



b) Estimate the asymptotic number of terms in expanded form of the  $n$ -variate polynomial  $\prod_{j=1}^n (X_j + 1)$ .

**PROBLEM 14 (4+4+2P):**

- Suppose  $\mathcal{A}$  is a randomized algorithm solving the decision problem  $L$  in time  $t(n)$  with one-sided error  $1 - \epsilon$  independent of  $n$ : On inputs  $\vec{x} \notin L$ ,  $\mathcal{A}$  always correctly reports **false**; but on inputs  $\vec{x} \in L$ ,  $\mathcal{A}$  might also report **false** with probability  $1 - \epsilon$ . Design and analyze an algorithm  $\mathcal{A}'$  that, by repeating  $\mathcal{A}$  an appropriate (which?) number of times, errs only with probability  $2^{-100n}$ .
- Now suppose randomized  $\mathcal{B}$  makes two-sided errors: reporting **false** on inputs  $\vec{x} \in L$  with probability  $\leq 1/2 - \epsilon$  and **true** on inputs  $\vec{x} \notin L$  with probability  $\leq 1/2 - \epsilon$  for some fixed  $\epsilon > 0$ . Design and analyze an algorithm  $\mathcal{B}'$  that repeatedly invokes  $\mathcal{B}$  to err with probability  $\leq 2^{-100n}$ .
- Search, select, study, and summarize two publications on causes and rates of so-called *soft errors*.